

# SOS and SDP Relaxation of Polynomial Optimization Problems

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.. - p.1/58

Purpose of this talk —  
Introduction to  
Sum Of Squares (SOS) and  
Semidefinite Programming (SDP)  
Relaxation of  
Polynomial Optimization Problems (POPs)

## Contents

1. Global optimization of nonconvex problems
2. Polynomial Optimization Problems (POPs)
3. Basic idea of SOS and SDP relaxations
4. SOS relaxation of POPs
5. SDP relaxation of POPs
6. Exploiting sparsity in SOS and SDP relaxations of POPs
7. Numerical results

.. - p.2/58

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.. - p.3/58

OP : Optimization problem in the  $n$ -dim. Euclidean space  $\mathbb{R}^n$   
min.  $f(x)$  sub.to  $x \in S \subseteq \mathbb{R}^n$ , where  $f : \mathbb{R}^n \rightarrow \mathbb{R}$ .

We want to approximate a global optimal solution  $x^*$ ;

$$x^* \in S \text{ and } f(x^*) \leq f(x) \text{ for all } x \in S$$

if it exists. But, impossible without any assumption.

Various assumptions

- continuity, differentiability, compactness, ...
- convexity  $\Rightarrow$  local opt. sol.  $\equiv$  global opt. sol.  
 $\Rightarrow$  local improvement leads to a global opt. sol.
- Powerful software for convex problems  $\ni$  LPs, SDPs, ...

.. - p.4/58

OP : Optimization problem in the  $n$ -dim. Euclidean space  $\mathbb{R}^n$   
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What can we do beyond convexity?

- We still need proper models and assumptions
  - Polynomial Optimization Problems (POPs) — this talk
- Main tool is SDP relaxation — this talk  
 Powerful in theory but expensive in practice
- Exploiting sparsity in large scale SDPs & POPs — this talk

.. p.5/58

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.. p.6/58

POP: min  $f_0(x)$  sub.to  $f_k(x) \geq 0$  or  $= 0$  ( $k = 1, \dots, m$ )

$x = (x_1, \dots, x_n) \in \mathbb{R}^n$  : a vector variable

$f_k(x)$  : a real-valued polynomial in  $x_1, \dots, x_n$  ( $k = 0, 1, \dots, m$ )

Example.  $n = 3$ ,  $x = (x_1, x_2, x_3) \in \mathbb{R}^3$  : a vector variable

$$\begin{aligned} \text{min} \quad & f_0(x) \equiv x_1^3 - 2x_1x_2^2 + x_1^2x_2x_3 - 4x_3^2 \\ \text{sub.to} \quad & f_1(x) \equiv -x_1^2 + 5x_2x_3 + 1 \geq 0, \\ & f_2(x) \equiv x_1^2 - 3x_1x_2x_3 + 2x_3 + 2 \geq 0, \\ & f_3(x) \equiv -x_1^2 - x_2^2 - x_3^2 + 1 \geq 0. \\ & x_1(x_1 - 1) = 0 \text{ (0-1 integer cond.)}, \\ & x_2 \geq 0, x_3 \geq 0, x_2x_3 = 0 \text{ (comp. cond.)}. \end{aligned}$$

- Various problems (including 0-1 integer programs)  $\Rightarrow$  POP
- POP serves as a unified theoretical model for global optimization in nonlinear and combinatorial optimization problems.

.. p.7/58

POP: min  $f_0(x)$  sub.to  $f_k(x) \geq 0$  or  $= 0$  ( $k = 1, \dots, m$ )

.. p.8/58

POP:  $\min f_0(\mathbf{x})$  sub.to  $f_k(\mathbf{x}) \geq 0$  or  $= 0$  ( $k = 1, \dots, m$ )

[1] Lasserre, "Global optimization with polynomials and the problems of moments", *SIAM J. on Optim.* (2001).

[2] Parrilo, "Semidefinite programming relaxations for semialgebraic problems", *Math. Prog.* (2003).

- **primal approach**  $\Rightarrow$  a sequence of SDP relaxations.
- **dual approach**  $\Rightarrow$  a sequence of SOS relaxations.

.. p.8/58

POP:  $\min f_0(\mathbf{x})$  sub.to  $f_k(\mathbf{x}) \geq 0$  or  $= 0$  ( $k = 1, \dots, m$ )

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- **primal approach**  $\Rightarrow$  a sequence of SDP relaxations.
- **dual approach**  $\Rightarrow$  a sequence of SOS relaxations.

Main features:

- (a) Lower bounds for the optimal value.
- (b) Convergence to global optimal solutions under assump.
- (c) Each relaxed problem can be solved as an SDP; its size  $\uparrow$  rapidly along "the sequence" as the size of POP  $\uparrow$ , the deg. of poly.  $\uparrow$ , and/or we require higher accuracy.
- (d) Expensive to solve large scale POPs in practice.  
 $\Rightarrow$  Exploiting Sparsity.

.. p.8/58

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Three ways of describing the SDP relaxation by Lasserre:

- (a) **Probability measure and its moments**
- (b) Linearization of polynomial SDPs
- (c) Sum of squares of polynomials

.. p.9/58

$\mu$  : a probability measure on  $\mathbb{R}^n$  . We assume  $n = 2$ . For every  $r = 0, 1, 2, \dots$  , define

$\mathbf{u}_r(\mathbf{x}) = (1, x_1, x_2, x_1^2, x_1x_2, x_2^2, x_1^3, \dots, x_2^r)^T$  : a column vector  
(all monomials with degree  $\leq r$ )

$\mathbf{M}_r(\mathbf{y}) = \int_{\mathbb{R}^2} \mathbf{u}_r(\mathbf{x}) \mathbf{u}_r(\mathbf{x})^T d\mu$   $\left( \begin{array}{l} \text{moment matrix, symmetric,} \\ \text{positive semidefinite} \end{array} \right)$

$y_{\alpha\beta} = \int_{\mathbb{R}^2} x_1^\alpha x_2^\beta d\mu = (\alpha, \beta)$ -element depending on  $\mu$ ,  $y_{00} = 1$

.. p.10/58

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(all monomials with degree  $\leq r$ )

$\mathbf{M}_r(\mathbf{y}) = \int_{\mathbb{R}^2} \mathbf{u}_r(\mathbf{x})\mathbf{u}_r(\mathbf{x})^T d\mu$  (moment matrix, symmetric, positive semidefinite)

$y_{\alpha\beta} = \int_{\mathbb{R}^2} x_1^\alpha x_2^\beta d\mu = (\alpha, \beta)$ -element depending on  $\mu$ ,  $y_{00} = 1$

Example with  $r = 2$ :

$$\mathbf{M}_r(\mathbf{y}) = \begin{pmatrix} y_{00} & y_{10} & y_{01} & y_{20} & y_{11} & y_{02} \\ y_{10} & y_{20} & y_{11} & y_{30} & y_{21} & y_{12} \\ y_{01} & y_{11} & y_{02} & y_{21} & y_{12} & y_{03} \\ y_{20} & y_{30} & y_{21} & y_{40} & y_{31} & y_{22} \\ y_{11} & y_{21} & y_{12} & y_{31} & y_{22} & y_{13} \\ y_{02} & y_{12} & y_{03} & y_{22} & y_{13} & y_{04} \end{pmatrix}, y_{00} = 1$$

$y_{21} = \int_{\mathbb{R}^2} x_1^2 x_2 d\mu$

$\mu$  : a probability measure on  $\mathbb{R}^n$  . We assume  $n = 2$ . For every  $r = 0, 1, 2, \dots$  , define

$\mathbf{u}_r(\mathbf{x}) = (1, x_1, x_2, x_1^2, x_1x_2, x_2^2, x_1^3, \dots, x_2^r)^T$  : a column vector  
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$\mathbf{M}_r(\mathbf{y}) = \int_{\mathbb{R}^2} \mathbf{u}_r(\mathbf{x})\mathbf{u}_r(\mathbf{x})^T d\mu$  (moment matrix, symmetric, positive semidefinite)

$y_{\alpha\beta} = \int_{\mathbb{R}^2} x_1^\alpha x_2^\beta d\mu = (\alpha, \beta)$ -element depending on  $\mu$ ,  $y_{00} = 1$

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(all monomials with degree  $\leq r$ )

$\mathbf{M}_r(\mathbf{y}) = \int_{\mathbb{R}^2} \mathbf{u}_r(\mathbf{x})\mathbf{u}_r(\mathbf{x})^T d\mu$  (moment matrix, symmetric, positive semidefinite)

$y_{\alpha\beta} = \int_{\mathbb{R}^2} x_1^\alpha x_2^\beta d\mu = (\alpha, \beta)$ -element depending on  $\mu$ ,  $y_{00} = 1$

$\mu$  : a probability measure on  $\mathbb{R}^2$   
 $\Downarrow$   
 $y_{00} = 1, \mathbf{M}_r(\mathbf{y}) \succeq \mathbf{O}$  (positive semidefinite) ( $r = 1, 2, \dots$ )

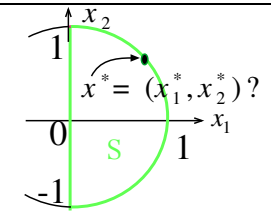
- We will use this necessary cond. with a finite  $r$  for  $\mu$  to be a probability measure in relaxation of a POP  $\Rightarrow$  next slide.

SDP relaxation (Lasserre '01) of a POP — an example

POP: min  $f_0(\mathbf{x}) = x_1^4 - 2x_1x_2$  opt. val.  $\zeta^*$  : unknown  
sub. to  $\mathbf{x} \in S \equiv \left\{ \mathbf{x} \in \mathbb{R}^2 : \begin{matrix} f_1(\mathbf{x}) = x_1 \geq 0 \\ f_2(\mathbf{x}) = 1 - x_1^2 - x_2^2 \geq 0 \end{matrix} \right\}$ .

$\Downarrow$   
min  $\int (x_1^4 x_2^0 - 2x_1 x_2) d\mu$   
sub. to  $\mu$  : a prob. meas. on  $S$ .

$\Downarrow y_{\alpha\beta} = \int_{\mathbb{R}^2} x_1^\alpha x_2^\beta d\mu$



### SDP relaxation (Lasserre '01) of a POP — an example

$$\text{POP: } \min_{\mathbf{x}} f_0(\mathbf{x}) = x_1^4 - 2x_1x_2 \quad \text{opt. val. } \zeta^* : \text{unknown}$$

$$\text{sub. to } \mathbf{x} \in S \equiv \left\{ \mathbf{x} \in \mathbb{R}^2 : \begin{array}{l} f_1(\mathbf{x}) = x_1 \geq 0 \\ f_2(\mathbf{x}) = 1 - x_1^2 - x_2^2 \geq 0 \end{array} \right\}.$$

⇕

$$\min \int (x_1^4 x_2^0 - 2x_1x_2) d\mu$$

$$\text{sub. to } \mu : \text{ a prob. meas. on } S.$$

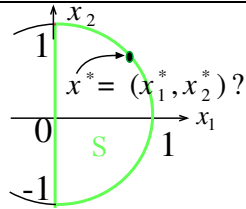
$$\Downarrow y_{\alpha\beta} = \int_{\mathbb{R}^2} x_1^\alpha x_2^\beta d\mu$$

$$\min y_{40} - 2y_{11} \Rightarrow \text{SDP relaxation, opt. val. } \zeta_r \leq \zeta^*$$

$$\text{sub. to "a certain moment cond. with a parameter } r$$

$$\text{for } \mu \text{ to be a probability measure on } S" \Rightarrow \text{next slide}$$

- $\zeta_r \leq \zeta_{r+1} \leq \zeta^*$ , and  $\zeta_r \rightarrow \zeta^*$  as  $r \rightarrow \infty$  under a moderate assumption that requires  $S$  is bounded (Lasserre '01).



### SDP relaxation (Lasserre '01) of a POP — an example

$$r = 2$$

$$\min_{\mu} y_{40} - 2y_{11} \text{ s.t. } \int_S \begin{pmatrix} 1 \\ x_1 \\ x_1^2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_1^2 \end{pmatrix}^T x_1 d\mu \succeq \mathbf{O}, \Leftrightarrow x_1 \geq 0$$

$$1 - x_1^2 - x_2^2 \geq 0 \Rightarrow \int_S \begin{pmatrix} 1 \\ x_1 \\ x_2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \end{pmatrix}^T (1 - x_1^2 - x_2^2) d\mu \succeq \mathbf{O},$$

$$\text{(moment matrix)} \int_S \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix}^T d\mu \succeq \mathbf{O}.$$

Here  $\mu$  denotes a probability measure on  $S$ .

↓

### SDP relaxation (Lasserre '01) of a POP — an example

- $\mu : \text{ a probability measure}$

$$y_{\alpha\beta} = \int_S x_1^\alpha x_2^\beta d\mu$$

$$f_1(\mathbf{x}) = x_1 \geq 0 \Rightarrow \int_S \begin{pmatrix} 1 \\ x_1 \\ x_1^2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_1^2 \end{pmatrix}^T x_1 d\mu \succeq \mathbf{O}$$

$$\Leftrightarrow \int_S \begin{pmatrix} x_1 & x_1^2 & x_1^3 \\ x_1^2 & x_1^3 & x_1^4 \\ x_1^3 & x_1^4 & x_1^5 \end{pmatrix} d\mu \succeq \mathbf{O} \Leftrightarrow \begin{pmatrix} y_{10} & y_{20} & y_{30} \\ y_{20} & y_{30} & y_{40} \\ y_{30} & y_{40} & y_{50} \end{pmatrix} \succeq \mathbf{O}$$

$$f_2(\mathbf{x}) = 1 - x_1^2 - x_2^2 \geq 0 \Rightarrow \int_S \begin{pmatrix} 1 \\ x_1 \\ x_2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \end{pmatrix}^T f_2(\mathbf{x}) d\mu \succeq \mathbf{O}$$

$$\Rightarrow \begin{pmatrix} 1 - y_{20} - y_{02} & y_{10} - y_{30} - y_{12} & y_{01} - y_{21} - y_{03} \\ y_{10} - y_{30} - y_{12} & y_{20} - y_{40} - y_{22} & y_{11} - y_{31} - y_{13} \\ y_{01} - y_{21} - y_{03} & y_{11} - y_{31} - y_{13} & y_{02} - y_{22} - y_{04} \end{pmatrix} \succeq \mathbf{O}$$

### SDP relaxation (Lasserre '01) of a POP — an example

- $\mu : \text{ a probability measure}$

$$y_{\alpha\beta} = \int_S x_1^\alpha x_2^\beta d\mu$$

$$\int_S \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix}^T d\mu \succeq \mathbf{O}$$

$$\Rightarrow \begin{pmatrix} 1 & y_{10} & y_{01} & y_{20} & y_{11} & y_{02} \\ y_{10} & y_{20} & y_{11} & y_{30} & y_{21} & y_{12} \\ y_{01} & y_{11} & y_{02} & y_{21} & y_{12} & y_{03} \\ y_{20} & y_{30} & y_{21} & y_{40} & y_{31} & y_{22} \\ y_{11} & y_{21} & y_{12} & y_{31} & y_{22} & y_{13} \\ y_{02} & y_{12} & y_{03} & y_{22} & y_{13} & y_{04} \end{pmatrix} \succeq \mathbf{O}$$

Consequently, we obtain

### SDP relaxation (Lasserre '01) of a POP — an example

$$\min y_{40} - 2y_{11} \text{ s.t. } \begin{pmatrix} y_{10} & y_{20} & y_{30} \\ y_{20} & y_{30} & y_{40} \\ y_{30} & y_{40} & y_{50} \end{pmatrix} \succeq \mathbf{O},$$

$$\begin{pmatrix} 1 - y_{20} - y_{02} & y_{10} - y_{30} - y_{12} & y_{01} - y_{21} - y_{03} \\ y_{10} - y_{30} - y_{12} & y_{20} - y_{40} - y_{22} & y_{11} - y_{31} - y_{13} \\ y_{01} - y_{21} - y_{03} & y_{11} - y_{31} - y_{13} & y_{02} - y_{22} - y_{04} \end{pmatrix} \succeq \mathbf{O},$$

$$\text{(moment matrix)} \begin{pmatrix} 1 & y_{10} & y_{01} & y_{20} & y_{11} & y_{02} \\ y_{10} & y_{20} & y_{11} & y_{30} & y_{21} & y_{12} \\ y_{01} & y_{11} & y_{02} & y_{21} & y_{12} & y_{03} \\ y_{20} & y_{30} & y_{21} & y_{40} & y_{31} & y_{22} \\ y_{11} & y_{21} & y_{12} & y_{31} & y_{22} & y_{13} \\ y_{02} & y_{12} & y_{03} & y_{22} & y_{13} & y_{04} \end{pmatrix} \succeq \mathbf{O}.$$

- We can apply SDP relaxation to general POPs in  $\mathbb{R}^n$ .
- Powerful in theory but very expensive in computation  
 $\Rightarrow$  Exploiting sparsity is crucial in practice.

— p.16/58

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- (b) **Linearization of polynomial SDPs**  
 $\rightarrow$  similar to (a), but easier to understand
- (c) Sum of squares of polynomials

— p.17/58

### SDP relaxation (Lasserre '01) of a POP — an example

$$\text{POP: } \min f_0(\mathbf{x}) = x_1^4 - 2x_1x_2 \quad \text{opt. val. } \zeta^* : \text{ unknown}$$

$$\text{sub. to } \mathbf{x} \in S \equiv \left\{ \mathbf{x} \in \mathbb{R}^2 : \begin{array}{l} f_2(\mathbf{x}) = x_1 \geq 0 \\ f_1(\mathbf{x}) = 1 - x_1^2 - x_2^2 \geq 0 \end{array} \right\}.$$

— p.18/58

### SDP relaxation (Lasserre '01) of a POP — an example

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$\Updownarrow$

### Polynomial SDP

$$\min x_1^4 - 2x_1x_2$$

$$\text{s.t. } \begin{pmatrix} 1 \\ x_1 \\ x_1^2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_1^2 \end{pmatrix}^T f_1(\mathbf{x}) \succeq \mathbf{O},$$

$$\begin{pmatrix} 1 \\ x_1 \\ x_2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \end{pmatrix}^T f_2(\mathbf{x}) \succeq \mathbf{O}, \quad \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix}^T \succeq \mathbf{O}.$$

- Expand poly.mat.inequalities.
- Linearize the problem by replacing  $x_1^\alpha x_2^\beta$  by  $y_{\alpha\beta} \Rightarrow$

— p.18/58

### SDP relaxation (Lasserre '01) of a POP — an example

$$\min x_1^4 - 2x_1x_2 \Rightarrow \min y_{40} - 2y_{11}$$

$$f_1(\mathbf{x}) = x_1 \geq 0 \Leftrightarrow \begin{pmatrix} 1 \\ x_1 \\ x_1^2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_1^2 \end{pmatrix}^T \succeq \mathbf{0}$$

$$\Leftrightarrow \begin{pmatrix} x_1 & x_1^2 & x_1^3 \\ x_1^2 & x_1^3 & x_1^4 \\ x_1^3 & x_1^4 & x_1^5 \end{pmatrix} \succeq \mathbf{0} \Rightarrow \begin{pmatrix} y_{10} & y_{20} & y_{30} \\ y_{20} & y_{30} & y_{40} \\ y_{30} & y_{40} & y_{50} \end{pmatrix} \succeq \mathbf{0}$$

— p.19/58

### SDP relaxation (Lasserre '01) of a POP — an example

$$f_2(\mathbf{x}) = 1 - x_1^2 - x_2^2 \geq 0 \Leftrightarrow \begin{pmatrix} 1 \\ x_1 \\ x_2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \end{pmatrix}^T \succeq \mathbf{0}$$

$$\Leftrightarrow \begin{pmatrix} 1 - x_1^2x_2^0 - x_1^0x_2^2 & x_1^1x_2^0 - x_1^3x_2^0 - x_1^1x_2^2 & x_1^0x_2^1 - x_1^2x_2^1 - x_1^0x_2^3 \\ x_1^1x_2^0 - x_1^3x_2^0 - x_1^1x_2^2 & x_1^2x_2^0 - x_1^4x_2^0 - x_1^2x_2^2 & x_1^1x_2^1 - x_1^3x_2^1 - x_1^1x_2^3 \\ x_1^0x_2^1 - x_1^2x_2^1 - x_1^0x_2^3 & x_1^1x_2^1 - x_1^3x_2^1 - x_1^1x_2^3 & x_1^0x_2^2 - x_1^2x_2^2 - x_1^0x_2^4 \end{pmatrix} \succeq \mathbf{0}$$

$$\Downarrow$$

$$\begin{pmatrix} 1 - y_{20} - y_{02} & y_{10} - y_{30} - y_{12} & y_{01} - y_{21} - y_{03} \\ y_{10} - y_{30} - y_{12} & y_{20} - y_{40} - y_{22} & y_{11} - y_{31} - y_{13} \\ y_{01} - y_{21} - y_{03} & y_{11} - y_{31} - y_{13} & y_{02} - y_{22} - y_{04} \end{pmatrix} \succeq \mathbf{0}$$

— p.20/58

### SDP relaxation (Lasserre '01) of a POP — an example

Similarly,

$$\begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix}^T \succeq \mathbf{0}$$

$$\Rightarrow \begin{pmatrix} 1 & y_{10} & y_{01} & y_{20} & y_{11} & y_{02} \\ y_{10} & y_{20} & y_{11} & y_{30} & y_{21} & y_{12} \\ y_{01} & y_{11} & y_{02} & y_{21} & y_{12} & y_{03} \\ y_{20} & y_{30} & y_{21} & y_{40} & y_{31} & y_{22} \\ y_{11} & y_{21} & y_{12} & y_{31} & y_{22} & y_{13} \\ y_{02} & y_{12} & y_{03} & y_{22} & y_{13} & y_{04} \end{pmatrix} \succeq \mathbf{0}$$

Consequently, we obtain

— p.21/58

### SDP relaxation (Lasserre '01) of a POP — an example

$$\min y_{40} - 2y_{11} \text{ s.t. } \begin{pmatrix} y_{10} & y_{20} & y_{30} \\ y_{20} & y_{30} & y_{40} \\ y_{30} & y_{40} & y_{50} \end{pmatrix} \succeq \mathbf{0},$$

$$\begin{pmatrix} 1 - y_{20} - y_{02} & y_{10} - y_{30} - y_{12} & y_{01} - y_{21} - y_{03} \\ y_{10} - y_{30} - y_{12} & y_{20} - y_{40} - y_{22} & y_{11} - y_{31} - y_{13} \\ y_{01} - y_{21} - y_{03} & y_{11} - y_{31} - y_{13} & y_{02} - y_{22} - y_{04} \end{pmatrix} \succeq \mathbf{0},$$

$$\begin{pmatrix} 1 & y_{10} & y_{01} & y_{20} & y_{11} & y_{02} \\ y_{10} & y_{20} & y_{11} & y_{30} & y_{21} & y_{12} \\ y_{01} & y_{11} & y_{02} & y_{21} & y_{12} & y_{03} \\ y_{20} & y_{30} & y_{21} & y_{40} & y_{31} & y_{22} \\ y_{11} & y_{21} & y_{12} & y_{31} & y_{22} & y_{13} \\ y_{02} & y_{12} & y_{03} & y_{22} & y_{13} & y_{04} \end{pmatrix} \succeq \mathbf{0}.$$

- If  $\mathbf{x}$  is a feasible sol. of Poly. SDP (or POP), then  $\mathbf{y} = (y_{\alpha\beta} = x_1^\alpha x_2^\beta)$  is a feasible sol. of SDP with the same obj. val.;  $x_1^4 - 2x_1x_2 = y_{40} - 2y_{11} \Rightarrow$  Relaxation

— p.22/58

## Contents

1. Global optimization of nonconvex problems
2. Polynomial Optimization Problems (POPs)
3. **Basic idea of SOS and SDP relaxations**

Three ways of describing the SDP relaxation by Lasserre:

- (a) Probability measure and its moments
- (b) Linearization of polynomial SDPs
- (c) **Sum of squares of polynomials**
  - dual approach
  - easy to understand in unconstrained case

— p.23/58

$f(\mathbf{x})$  : an SOS (Sum of Squares) polynomial

$$\begin{aligned} & \Downarrow \\ & \exists \text{ polynomials } g_1(\mathbf{x}), \dots, g_q(\mathbf{x}); f(\mathbf{x}) = \sum_{j=1}^q g_j(\mathbf{x})^2. \end{aligned}$$

$\mathcal{N}$  : the set of nonnegative polynomials in  $\mathbf{x} \in \mathbb{R}^n$ .

$\text{SOS}_*$  : the set of SOS. Obviously,  $\text{SOS}_* \subset \mathcal{N}$ .

$\text{SOS}_r = \{f \in \text{SOS}_* : \deg f \leq 2r\}$  : SOSs w. degree at most  $2r$ .

$$n = 2. f(x_1, x_2) = (x_1^2 - 2x_2 + 1)^2 + (3x_1x_2 + x_2 - 4)^2 \in \text{SOS}_2.$$

$$n = 2. f(x_1, x_2) = (x_1x_2 - 1)^2 + x_1^2 \in \text{SOS}_2.$$

- In theory,  $\text{SOS}_*$  (SOS)  $\subset \mathcal{N}$ .  $\text{SOS}_* \neq \mathcal{N}$  in general.
- If  $n = 1$ ,  $\text{SOS}_* = \mathcal{N}$ .  $\{f \in \mathcal{N} : \deg f \leq 2\} \equiv \text{SOS}_1$ .
- In practice,  $f(\mathbf{x}) \in \mathcal{N} \setminus \text{SOS}_*$  is rare.
- So we replace  $\mathcal{N}$  by  $\text{SOS}_* \implies$  SOS Relaxations.

— p.24/58

Representation of

$$\text{SOS}_r \equiv \left\{ \sum_{j=1}^q g_j(\mathbf{x})^2 : q \geq 1, g_j(\mathbf{x}) \text{ is a poly. of deg } \leq r \right\} \subset \text{SOS}_*.$$

$\forall$  poly.  $g(\mathbf{x})$  of deg  $\leq r$ ,  $\exists \mathbf{a} \in \mathbb{R}^{d(r)}$ ;  $g(\mathbf{x}) = \mathbf{a}^T \mathbf{u}_r(\mathbf{x})$ , where

$$\mathbf{u}_r(\mathbf{x}) = (1, x_1, \dots, x_n, x_1^2, x_1x_2, x_1x_3, \dots, x_n^2, \dots, x_1^r, \dots, x_n^r)^T$$

(a column vector of a basis for polynomials of degree  $\leq r$ ),

$$d(r) = \binom{n+r}{r} : \text{the dimension of } \mathbf{u}_r(\mathbf{x}).$$

$\Downarrow$

$$\begin{aligned} \text{SOS}_r &= \left\{ \sum_{j=1}^q (\mathbf{a}_j^T \mathbf{u}_r(\mathbf{x}))^2 : q \geq 1, \mathbf{a}_j \in \mathbb{R}^{d(r)} \right\} \\ &= \left\{ \mathbf{u}_r(\mathbf{x})^T \left( \sum_{j=1}^q \mathbf{a}_j \mathbf{a}_j^T \right) \mathbf{u}_r(\mathbf{x}) : q \geq 1, \mathbf{a}_j \in \mathbb{R}^{d(r)} \right\} \\ &= \left\{ \mathbf{u}_r(\mathbf{x})^T \mathbf{V} \mathbf{u}_r(\mathbf{x}) : \mathbf{V} \succeq \mathbf{O} \right\}. \end{aligned}$$

— p.25/58

Example.  $n = 1$ , SOS polynomials of degree  $\leq 3$  in  $x \in \mathbb{R}$ .

$$\begin{aligned} \text{SOS}_3 &\equiv \left\{ \sum_{j=1}^q g_j(x)^2 : q \geq 1, g_j(x) \text{ is a poly. of degree } \leq 3 \right\} \\ &= \left\{ \begin{pmatrix} 1 \\ x \\ x^2 \\ x^3 \end{pmatrix}^T \mathbf{V} \begin{pmatrix} 1 \\ x \\ x^2 \\ x^3 \end{pmatrix} : \mathbf{V} \text{ is } 4 \times 4 \text{ psd matrix} \right\} \end{aligned}$$

— p.26/58



Example.  $n = 2$ , SOS polynomials of degree  $\leq 2$  in  $x=(x_1, x_2)$ .

$$\text{sos}_2 \equiv \left\{ \sum_{j=1}^q g_j(\mathbf{x})^2 : q \geq 1, g_j(\mathbf{x}) \text{ is a poly. of degree } \leq 2 \right\}$$

$$= \left\{ \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix}^T \mathbf{V} \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix} : \mathbf{V} \text{ is a } 6 \times 6 \text{ psd matrix} \right\}$$

$$f(\mathbf{x}) = -x_1 + 2x_2 + 3x_1^2 - 5x_1^2x_2^2 + 7x_2^4 \quad \zeta = 3.1: \text{fixed}$$

$$f(\mathbf{x}) - \zeta \geq 0 \ (\forall \mathbf{x}) \implies \text{SOS relaxation}$$

$$f(\mathbf{x}) - \zeta \in \text{sos}_2 \ (\text{SOS of poly. of degree } \leq 2)$$



$$\exists \mathbf{V} \in \mathbb{S}^6; \quad f(\mathbf{x}) - \zeta =$$

Sum of Squares

$$\begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix}^T \begin{pmatrix} V_{11} & V_{12} & V_{13} & V_{14} & V_{15} & V_{16} \\ V_{12} & V_{22} & V_{23} & V_{24} & V_{25} & V_{26} \\ V_{13} & V_{23} & V_{33} & V_{34} & V_{35} & V_{36} \\ V_{14} & V_{24} & V_{34} & V_{44} & V_{45} & V_{46} \\ V_{15} & V_{25} & V_{35} & V_{45} & V_{55} & V_{56} \\ V_{16} & V_{26} & V_{36} & V_{46} & V_{56} & V_{66} \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix}$$

$$(\forall (x_1, x_2)^T \in \mathbb{R}^n), \quad \mathbf{V} \succeq \mathbf{O}$$

$\Updownarrow$  Compare the coef. of  $1, x_1, x_2, x_1^2, x_1x_2, x_2^2$  on both sides of  $=$

LMI:  $\exists \mathbf{V} \in \mathbb{S}^6?$

$$-\zeta = V_{11}, \quad -1 = 2V_{12}, \quad 2 = 2V_{13}, \quad 3 = 2V_{14} + V_{22},$$

$$-5 = 2V_{46} + V_{55}, \quad 7 = V_{66}, \quad \text{all others } 0 = \dots,$$

$$\mathbf{V} \succeq \mathbf{O}$$

In general, each equality constraint is a linear equation in  $\zeta$  and  $\mathbf{V}$ .

Unconstrained minimization

$$\mathcal{P}: \min f(\mathbf{x}) = -x_1 + 2x_2 + 3x_1^2 - 5x_1^2x_2^2 + 7x_2^4$$

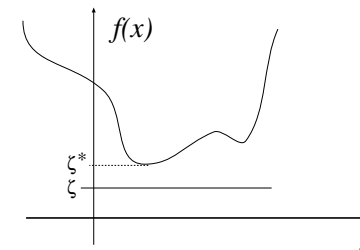
$\Updownarrow$  (a special case of Lagrangian dual)

$$\mathcal{P}': \max \zeta \text{ s.t. } f(\mathbf{x}) - \zeta \geq 0 \ (\forall \mathbf{x} \in \mathbb{R}^n)$$



$$f(\mathbf{x}) - \zeta \in \mathcal{N} \ (\text{the nonnegative polynomials})$$

Here  $\mathbf{x}$  is a parameter (index) describing inequality constraints.



### Unconstrained minimization

$$\mathcal{P}: \min f(\mathbf{x}) = -x_1 + 2x_2 + 3x_1^2 - 5x_1^2x_2^2 + 7x_2^4$$

⇕ (a special case of Lagrangian dual)

$$\mathcal{P}': \max \zeta \quad \text{s.t.} \quad f(\mathbf{x}) - \zeta \geq 0 \quad (\forall \mathbf{x} \in \mathbb{R}^n)$$

$$\Downarrow$$

$$f(\mathbf{x}) - \zeta \in \mathcal{N} \text{ (the nonnegative polynomials)}$$

Here  $x$  is a parameter (index) describing inequality constraints.

### Unconstrained minimization

$$\mathcal{P}: \min f(\mathbf{x}) = -x_1 + 2x_2 + 3x_1^2 - 5x_1^2x_2^2 + 7x_2^4$$

⇕ (a special case of Lagrangian dual)

$$\mathcal{P}': \max \zeta \quad \text{s.t.} \quad f(\mathbf{x}) - \zeta \geq 0 \quad (\forall \mathbf{x} \in \mathbb{R}^n)$$

$$\Downarrow$$

$$f(\mathbf{x}) - \zeta \in \mathcal{N} \text{ (the nonnegative polynomials)}$$

Here  $x$  is a parameter (index) describing inequality constraints.

⇓ a subproblem of  $\mathcal{P}'$  (or a relaxation of  $\mathcal{P}$ )

$$\mathcal{P}'': \max \zeta \quad \text{sub.to} \quad f(\mathbf{x}) - \zeta \in \text{SOS}_2 \text{ (SOS of poly. of degree } \leq 2)$$

⇓

$$\max \zeta \quad \text{s.t.} \quad f(\mathbf{x}) - \zeta =$$

Sum of Squares

$$\begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix}^T \begin{pmatrix} V_{11} & V_{12} & V_{13} & V_{14} & V_{15} & V_{16} \\ V_{12} & V_{22} & V_{23} & V_{24} & V_{25} & V_{26} \\ V_{13} & V_{23} & V_{33} & V_{34} & V_{35} & V_{36} \\ V_{14} & V_{24} & V_{34} & V_{44} & V_{45} & V_{46} \\ V_{15} & V_{25} & V_{35} & V_{45} & V_{55} & V_{56} \\ V_{16} & V_{26} & V_{36} & V_{46} & V_{56} & V_{66} \end{pmatrix} \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_1^2 \\ x_1x_2 \\ x_2^2 \end{pmatrix}$$

$$(\forall (x_1, x_2)^T \in \mathbb{R}^n), \quad \mathbf{V} \succeq \mathbf{O}$$

⇕ Compare the coef. of  $1, x_1, x_2, x_1^2, x_1x_2, x_2^2$  on both side of =

### SDP (Semidefinite Program)

$$\max \zeta \quad \text{s.t.} \quad -\zeta = V_{11}, \quad -1 = 2V_{12}, \quad 2 = 2V_{13}, \quad 3 = 2V_{14} + V_{22},$$

$$-5 = 2V_{46} + V_{55}, \quad 7 = V_{66}, \quad \text{all others } 0 = \dots,$$

$$\mathbf{V} \succeq \mathbf{O}$$

In general, each equality constraint is a linear equation in  $\zeta$  and  $\mathbf{V}$ .

### SDP (Semidefinite Program)

$$\begin{aligned} \max \quad & \zeta \text{ s.t.} \quad -\zeta = V_{11}, \quad -1 = 2V_{12}, \quad 2 = 2V_{13}, \quad 3 = 2V_{14} + V_{22}, \\ & -5 = 2V_{46} + V_{55}, \quad 7 = V_{66}, \quad \text{all others } 0 = \dots, \\ & \mathbf{V} \succeq \mathbf{O} \end{aligned}$$

In general, each equality constraint is a linear equation in  $\zeta$  and  $V$ .

- SOS relaxation of general constrained POPs  
⇒ Next section
- We will use a generalized Lagrangian dual.

.. p.33/58

### Contents

1. Global optimization of nonconvex problems
2. Polynomial Optimization Problems (POPs)
3. Basic idea of SOS and SDP relaxations
4. **SOS relaxation of POPs**
5. SDP relaxation of POPs
6. Exploiting sparsity in SOS and SDP relaxations of POPs
7. Numerical results

.. p.34/58

$$\text{POP: } \zeta^* = \min. f_0(\mathbf{x}) \text{ s.t. } f_k(\mathbf{x}) \geq 0 \ (k = 1, \dots, m)$$

⇓

A sequence of SOS relaxation problems  $\text{SOS}^r$  depending on the relaxation order  $r = 1, 2, \dots$ , which determines the degree of polynomials used;

- (a) Under a moderate assumption, opt. sol. of  $\text{SOS}^r \rightarrow$  opt sol. of POP as  $r \rightarrow \infty$  (Lasserre 2001).
- (b)  $r = \lceil \text{"the max. deg. of poly. in POP"}/2 \rceil + 0 \sim 3$  is usually large enough to attain opt sol. of POP in practice.
- (c) Such an  $r$  is **unknown** in theory except  $\exists$  special cases.
- (d)  $\text{SOS}^r$  can be converted into an  $\text{SDP}^r$  and solved by the primal-dual interior-point method.
- (e) The size of  $\text{SDP}^r$  **increases** as  $r \rightarrow \infty$ .

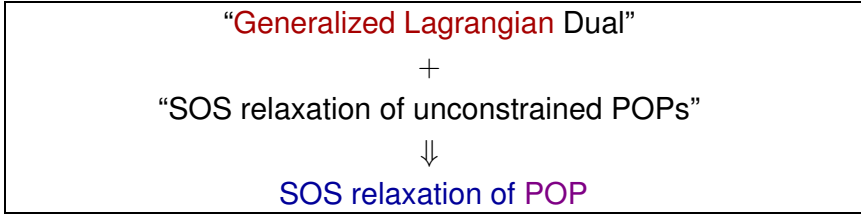
.. p.35/58

$$\text{POP: } \zeta^* = \min. f_0(\mathbf{x}) \text{ s.t. } f_k(\mathbf{x}) \geq 0 \ (k = 1, \dots, m)$$

.. p.36/58

$$\text{POP: } \zeta^* = \min. f_0(\mathbf{x}) \text{ s.t. } f_k(\mathbf{x}) \geq 0 \ (k = 1, \dots, m)$$

● Rough sketch of SOS relaxation of POP



$$\text{POP: } \zeta^* = \min. f_0(\mathbf{x}) \text{ s.t. } f_k(\mathbf{x}) \geq 0 \ (k = 1, \dots, m)$$

$$\text{POP: } \zeta^* = \min. f_0(\mathbf{x}) \text{ s.t. } f_k(\mathbf{x}) \geq 0 \ (k = 1, \dots, m)$$

↓

**Generalized Lagrangian function**

$$L(\mathbf{x}, \lambda_1, \dots, \lambda_m) = f_0(\mathbf{x}) - \sum_{k=1}^m \lambda_k(\mathbf{x}) f_k(\mathbf{x})$$

for  $\forall \mathbf{x} \in \mathbb{R}^n, \forall \lambda_k \in \text{SOS}_*$

If  $\mathbb{R} \ni \lambda_k \geq 0$  then **L** is the standard Lagrangian function.

**Generalized Lagrangian Dual**

$$\max_{\lambda_1 \in \text{SOS}_*, \dots, \lambda_m \in \text{SOS}_*} \min_{\mathbf{x} \in \mathbb{R}^n} L(\mathbf{x}, \lambda_1, \dots, \lambda_m)$$

$$\Leftrightarrow \max \eta \text{ s.t. } L(\mathbf{x}, \lambda_1, \dots, \lambda_m) - \eta \geq 0 \ (\forall \mathbf{x} \in \mathbb{R}^n)$$

**Generalized Lagrangian Dual**

$$\max \eta \text{ s.t. } L(\mathbf{x}, \lambda_1, \dots, \lambda_m) - \eta \geq 0 \ (\forall \mathbf{x} \in \mathbb{R}^n),$$

$$\lambda_1 \in \text{SOS}_*, \dots, \lambda_m \in \text{SOS}_*$$

$\mathbf{x}$  is not a variable but the index for infinite no. of inequalities.

$$\text{POP: } \zeta^* = \min. f_0(\mathbf{x}) \text{ s.t. } f_k(\mathbf{x}) \geq 0 \ (k = 1, \dots, m)$$

**POP:**  $\zeta^* = \min. f_0(\mathbf{x})$  s.t.  $f_k(\mathbf{x}) \geq 0$  ( $k = 1, \dots, m$ )

**Generalized Lagrangian function**

$$\Downarrow \quad L(\mathbf{x}, \lambda_1, \dots, \lambda_m) = f_0(\mathbf{x}) - \sum_{k=1}^m \lambda_k(\mathbf{x}) f_k(\mathbf{x})$$

for  $\forall \mathbf{x} \in \mathbb{R}^n, \forall \lambda_k \in \mathbf{SOS}_*$

**Generalized Lagrangian Dual**

$$\max \eta \text{ s.t. } L(\mathbf{x}, \lambda_1, \dots, \lambda_m) - \eta \geq 0 \quad (\forall \mathbf{x} \in \mathbb{R}^n),$$

$$\lambda_1 \in \mathbf{SOS}_*, \dots, \lambda_m \in \mathbf{SOS}_*$$

$\Downarrow$  **SOS relaxation**

$$\mathbf{SOS}^r: \eta^r = \max \eta \text{ s.t. } L(\mathbf{x}, \lambda_1, \dots, \lambda_m) - \eta \in \mathbf{SOS}_r \quad (\forall \mathbf{x} \in \mathbb{R}^n),$$

$$\lambda_1 \in \mathbf{SOS}_{r_1}, \dots, \lambda_m \in \mathbf{SOS}_{r_m}$$

- $\mathbf{SOS}_{r_k}$ : the set of sum of square polynomials with  $\text{deg.} \leq r_k$
- $\lceil \max\{\text{deg}(f_k) : k = 0, \dots, m\}/2 \rceil \leq r$ : the relaxation order
- $r_k = r - \lceil \text{deg}(f_k)/2 \rceil$  ( $k = 1, \dots, m$ ); chosen to balance the degrees of all the terms  $\lambda_k(\mathbf{x}) f_k(\mathbf{x})$  ( $k = 1, \dots, m$ ).
- $\eta^r \leq \eta^{r+1}$ ,  $\eta^r \rightarrow \zeta^*$  as  $r \rightarrow \infty$  under a moderate assumption which requires the feasible region of **POP** is compact.

... p.38/58

## Contents

1. Global optimization of nonconvex problems
2. Polynomial Optimization Problems (POPs)
3. Basic idea of SOS and SDP relaxations
4. SOS relaxation of POPs
5. **SDP relaxation of POPs**
6. Exploiting sparsity in SOS and SDP relaxations of POPs
7. Numerical results

... p.39/58

**POP:**  $\zeta^* = \min. f_0(\mathbf{x})$  s.t.  $f_k(\mathbf{x}) \geq 0$  ( $k = 1, \dots, m$ )

- the relax. order  $r \geq \lceil \max\{\text{deg}(f_k) : k = 0, \dots, m\}/2 \rceil$

$$\Updownarrow \quad r_k = r - \lceil \text{deg}(f_k)/2 \rceil \quad (k = 1, \dots, m)$$

$$\mathbf{u}_r(\mathbf{x}) = (1, x_1, x_2, x_1^2, x_1x_2, x_2^2, x_1^3, \dots, x_2^r)^T$$

(all monomials with degree  $\leq r$ )

**Poly.SDP:**  $\min f_0(\mathbf{x})$  s.t.  $\mathbf{u}_{r_k}(\mathbf{x}) \mathbf{u}_{r_k}(\mathbf{x})^T f_k(\mathbf{x}) \succeq \mathbf{O}$  ( $\forall k$ )

$$\mathbf{u}_r(\mathbf{x}) \mathbf{u}_r(\mathbf{x})^T \succeq \mathbf{O}.$$

Linearize:

- $$\Downarrow$$
- Expand all the polynomial inequalities
  - Replace  $x_1^\alpha x_2^\beta \dots x_n^\gamma$  by a single variable  $y_{\alpha\beta\dots\gamma}$

**Linear SDP<sup>r</sup>:**  $\zeta^r = \min.$  a linear function in  $y_{\alpha\beta\dots\gamma}$

s.t. linear matrix inequalities in  $y_{\alpha\beta\dots\gamma}$

- $\zeta^r \leq \zeta^{r+1}$ ,  $\zeta^r \rightarrow \zeta^*$  as  $r \rightarrow \infty$  under a moderate assumption which requires the feasible region of **POP** is compact.

... p.40/58

## Example

**POP:**  $\min. f_0(\mathbf{x}) = x_1^3 - 2x_2^2$  s.t.  $f_1(\mathbf{x}) = 1 - x_1^2 - x_2^2 \geq 0$ .

$$\Updownarrow \quad \begin{aligned} & \bullet r = 2 \geq \lceil \max\{\text{deg}(f_k) : k = 0, 1\}/2 \rceil = \lceil 3/2 \rceil = 2. \\ & \bullet r_1 = r - \lceil \text{deg}(f_1)/2 \rceil = 2 - \lceil 2/2 \rceil = 1. \\ & \bullet \mathbf{u}_1(\mathbf{x}) = (1, x_1, x_2)^T, \mathbf{u}_2(\mathbf{x}) = (1, x_1, x_2, x_1^2, x_1x_2, x_2^2)^T. \end{aligned}$$

**Poly.SDP:**  $\min x_1^3 - 2x_2^2$

s.t.  $\mathbf{u}_1(\mathbf{x}) \mathbf{u}_1(\mathbf{x})^T f_1(\mathbf{x}) \succeq \mathbf{O}$

$$\mathbf{u}_2(\mathbf{x}) \mathbf{u}_2(\mathbf{x})^T \succeq \mathbf{O}.$$

$\Downarrow$  Expand matrix inequalities

... p.41/58

**Poly.SDP:**  $\min x_1^3 - 2x_2^2$  s.t.

$$\begin{pmatrix} 1 - x_1^2 - x_2^2 & x_1 - x_1^3 - x_1x_2^2 & x_2 - x_1^2x_2 - x_2^3 \\ x_1 - x_1^3 - x_1x_2^2 & x_1^2 - x_1^4 - x_1^2x_2^2 & x_1x_2 - x_1^3x_2 - x_1x_2^3 \\ x_2 - x_1^2x_2 - x_2^3 & x_1x_2 - x_1^3x_2 - x_1x_2^3 & x_2^2 - x_1^2x_2^2 - x_2^4 \end{pmatrix} \succeq \mathbf{O},$$

$$\begin{pmatrix} 1 & x_1 & x_2 & x_1^2 & x_1x_2 & x_2^2 \\ x_1 & x_1^2 & x_1x_2 & x_1^3 & x_1^2x_2 & x_1x_2^2 \\ x_2 & x_1x_2 & x_2^2 & x_1^2x_2 & x_1x_2^2 & x_2^3 \\ x_1^2 & x_1^3 & x_1^2x_2 & x_1^4 & x_1^3x_2 & x_1^2x_2^2 \\ x_1x_2 & x_1^2x_2 & x_1x_2^2 & x_1^3x_2 & x_1^2x_2^2 & x_1x_2^3 \\ x_2^2 & x_1x_2^2 & x_2^3 & x_1^2x_2^2 & x_1x_2^3 & x_2^4 \end{pmatrix} \succeq \mathbf{O}.$$

⇓ Linearize by replacing  $x_1^\alpha x_2^\beta$  by a single variable  $y_{\alpha\beta}$

... p.42/58

**SDP<sup>2</sup>:**  $\min y_{30} - 2y_{02}$  s.t.

$$\begin{pmatrix} 1 - y_{20} - y_{02} & y_{10} - y_{30} - y_{12} & y_{01} - y_{21} - y_{03} \\ y_{10} - y_{30} - y_{12} & y_{20} - y_{40} - y_{22} & y_{11} - y_{31} - y_{13} \\ y_{01} - y_{21} - y_{03} & y_{11} - y_{31} - y_{13} & y_{02} - y_{22} - y_{04} \end{pmatrix} \succeq \mathbf{O},$$

$$\begin{pmatrix} 1 & y_{10} & y_{01} & y_{20} & y_{11} & y_{02} \\ y_{10} & y_{20} & y_{11} & y_{30} & y_{21} & y_{12} \\ y_{01} & y_{11} & y_{02} & y_{21} & y_{12} & y_{03} \\ y_{20} & y_{30} & y_{21} & y_{40} & y_{31} & y_{22} \\ y_{11} & y_{21} & y_{12} & y_{31} & y_{22} & y_{13} \\ y_{02} & y_{12} & y_{03} & y_{22} & y_{13} & y_{04} \end{pmatrix} \succeq \mathbf{O}.$$

... p.43/58

## Duality

**POP:**  $\zeta^* = \min. f_0(\mathbf{x})$  s.t.  $f_k(\mathbf{x}) \geq 0$  ( $k = 1, \dots, m$ )

⇓

- the relax. order  $r \geq \lceil \max\{\deg(f_k) : k = 0, \dots, m\}/2 \rceil$
- $r_k = r - \lceil \deg(f_k)/2 \rceil$  ( $k = 1, \dots, m$ )

**SOS<sup>r</sup>:**  $\eta^r = \max \eta$  s.t.  $L(\mathbf{x}, \lambda_1, \dots, \lambda_m) - \eta \in \text{SOS}_r$  ( $\forall \mathbf{x} \in \mathbb{R}^n$ ),  
 $\lambda_1 \in \text{SOS}_{r_1}, \dots, \lambda_m \in \text{SOS}_{r_m}$

⇓ Convert

**SDP<sub>d</sub><sup>r</sup>:**  $\eta^r = \max.$  a linear objective function  
s.t. linear matrix inequalities.

⇕ primal and dual to each other

**SDP<sup>r</sup>:**  $\zeta^r = \min.$  a linear function in  $y_{\alpha\beta\dots\gamma}$   
s.t. linear matrix inequalities in  $y_{\alpha\beta\dots\gamma}$

- $\eta^r \leq \zeta^r$ .
- $\eta^r = \zeta^r$  if both **SDP<sub>d</sub><sup>r</sup>** and **SDP<sup>r</sup>** have int. feasible solutions.

... p.44/58

## Contents

1. Global optimization of nonconvex problems
2. Polynomial Optimization Problems (POPs)
3. Basic idea of SOS and SDP relaxations
4. SOS relaxation of POPs
5. SDP relaxation of POPs
6. Exploiting sparsity in SOS and SDP relaxations of POPs
7. Numerical results

... p.45/58

### A sparse POP example, alkyl.gms from globalib

$$\begin{aligned} \min f_0 &= -6.3x_5x_8 + 5.04x_2 + 0.35x_3 + x_4 + 3.36x_6 \text{ sub.to} \\ f_1 &= -0.820x_2 + x_5 - 0.820x_6 = 0 \quad f_2 = -x_2x_9 + 10x_3 + x_6 = 0 \\ f_3 &= 0.98x_4 - x_7(0.01x_5x_{10} + x_4) = 0, \quad \text{lbd}_i \leq x_i \leq \text{ubd}_i \\ f_4 &= x_5x_{12} - x_2(1.12 + 0.132x_9 - 0.0067x_9^2) = 0 \\ f_5 &= x_8x_{13} - 0.01x_9(1.098 - 0.038x_9) - 0.325x_7 - 0.574 = 0 \\ f_6 &= x_{10}x_{14} + 22.2x_{11} - 35.82 = 0 \quad f_7 = x_1x_{11} - 3x_8 - 1.33 = 0 \end{aligned}$$

- $n = 14$  variables. polynomials with  $\text{deg} \leq 3$ .
- $\forall f_k$  involves at most 6 variables.

...p.46/58

### A sparse POP example, alkyl.gms from globalib

$$\begin{aligned} \min f_0 &= -6.3x_5x_8 + 5.04x_2 + 0.35x_3 + x_4 + 3.36x_6 \text{ sub.to} \\ f_1 &= -0.820x_2 + x_5 - 0.820x_6 = 0 \quad f_2 = -x_2x_9 + 10x_3 + x_6 = 0 \\ f_3 &= 0.98x_4 - x_7(0.01x_5x_{10} + x_4) = 0, \quad \text{lbd}_i \leq x_i \leq \text{ubd}_i \\ f_4 &= x_5x_{12} - x_2(1.12 + 0.132x_9 - 0.0067x_9^2) = 0 \\ f_5 &= x_8x_{13} - 0.01x_9(1.098 - 0.038x_9) - 0.325x_7 - 0.574 = 0 \\ f_6 &= x_{10}x_{14} + 22.2x_{11} - 35.82 = 0 \quad f_7 = x_1x_{11} - 3x_8 - 1.33 = 0 \end{aligned}$$

- $n = 14$  variables. polynomials with  $\text{deg} \leq 3$ .
- $\forall f_k$  involves at most 6 variables.

Let  $F_k = \{i : x_i \text{ is involved } f_k\}$ ;  $F_1 = \{2, 5, 6\}$ ,  $F_2 = \{2, 3, 6, 9\}$ .  
Sparsity pattern matrix  $R \in \mathbb{S}^n$  (symbolic) & graph  $G(N, E)$ .

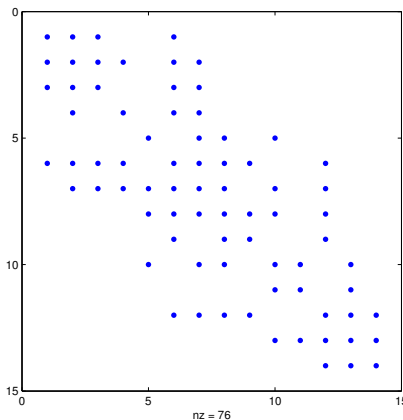
$$R_{ij} = \begin{cases} \star & \text{if } x_i, x_j \in \text{a monomial of } f_0 \text{ or } \{i, j\} \subset F_k \exists k \geq 1, \\ 0 & \text{otherwise.} \end{cases}$$

$$N = \{1, 2, \dots, n\}, \quad E = \{(i, j) \in N \times N : i \neq j \text{ and } R_{ij} = \star\}.$$

eg.,  $R_{25} = R_{52} = R_{58} = R_{85} = \star$ ,  $R_{28} = R_{82} = 0$

...p.46/58

The  $14 \times 14$  sparsity pattern matrix  $R$  with simultaneous row and column reordering (Reverse Cuthill-McKee ordering)



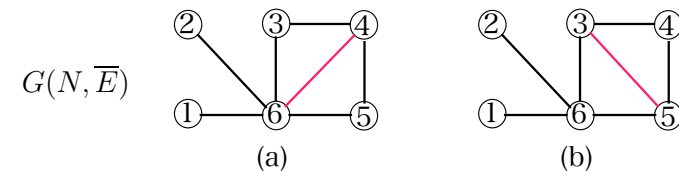
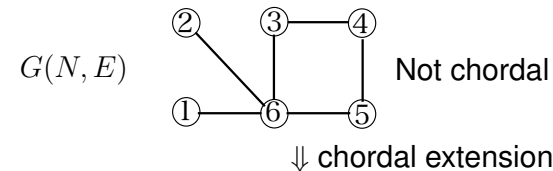
#### Structured sparsity

- Sparse (symbolic) Cholesky factorization
- Also, characterized by a sparse chordal graph structure

...p.47/58

$G(N, E)$  : a graph,  $N = \{1, \dots, n\}$  (nodes),  $E \subset N \times N$  (edges)

chordal  $\Leftrightarrow \forall$  cycle with more than 3 edges has a chord



$\{1, 6\}, \{2, 6\}, \{3, 4, 6\},$   
 $\{4, 5, 6\}$

$\{1, 6\}, \{2, 6\}, \{3, 5, 6\},$   
 $\{3, 4, 5\}$

Maximal cliques (node sets of maximal complete subgraphs)

...p.48/58

**POP:**  $\zeta^* = \min. f_0(\mathbf{x})$  s.t.  $f_k(\mathbf{x}) \geq 0$  ( $k = 1, \dots, m$ )

Let  $F_k = \{i : x_i \text{ is involved } f_k\}$ .

Sparsity pattern matrix  $\mathbf{R} \in \mathbb{S}^n$  (symbolic) & graph  $G(N, E)$ .

$$R_{ij} = \begin{cases} \star & \text{if } x_i, x_j \in \text{a monomial of } f_0 \text{ or } \{i, j\} \subset F_k \exists k \geq 1, \\ 0 & \text{otherwise.} \end{cases}$$

$$N = \{1, 2, \dots, n\}, E = \{(i, j) \in N \times N : i \neq j \text{ and } R_{ij} = \star\}.$$

-- p.49/58

**POP:**  $\zeta^* = \min. f_0(\mathbf{x})$  s.t.  $f_k(\mathbf{x}) \geq 0$  ( $k = 1, \dots, m$ )

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$$N = \{1, 2, \dots, n\}, E = \{(i, j) \in N \times N : i \neq j \text{ and } R_{ij} = \star\}.$$

- $G(N, E)$  is chordal iff  $\exists$  simultaneous row and column ordering of  $\mathbf{R}$  that allows a Cholesky factorization of  $\mathbf{R}$  with no fill-in.
- If  $G(N, E)$  is not chordal, then there are fill-ins in its Cholesky factorization, which are corresponding to the edges that are added in its chordal extension  $G(N, \overline{E})$ .
- The chordal extension is not unique; each chordal extension is corresponding a Cholesky factorization with  $\exists$  simultaneous row and column ordering of  $\mathbf{R}$

-- p.49/58

**POP:**  $\zeta^* = \min. f_0(\mathbf{x})$  s.t.  $f_k(\mathbf{x}) \geq 0$  ( $k = 1, \dots, m$ )

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-- p.50/58

**POP:**  $\zeta^* = \min. f_0(\mathbf{x})$  s.t.  $f_k(\mathbf{x}) \geq 0$  ( $k = 1, \dots, m$ )

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$$N = \{1, 2, \dots, n\}, E = \{(i, j) \in N \times N : i \neq j \text{ and } R_{ij} = \star\}.$$

$\Gamma$  : the set of max. cliques of a chordal ext.  $G(N, \overline{E})$  of  $G(N, E)$ .

Then  $F_k \subset \exists C \in \Gamma$ . Let  $F_k \subset C_k \in \Gamma$  ( $k = 1, 2, \dots, m$ ). Let

the relax. order  $r \geq \lceil \max\{\deg(f_k) : k = 0, \dots, m\}/2 \rceil$ ,

$$r_k = r - \lceil \deg(f_k)/2 \rceil \quad (k = 1, \dots, m),$$

$\text{SOS}_q(C)$  = the sum of square poly. with  $\deg. \leq q$  in  $x_i$  ( $i \in C$ ).

**Sparse SOS relaxation  $\text{SOS}^r$ :**

$$\max \eta \text{ s.t. } L(\mathbf{x}, \lambda_1, \dots, \lambda_m) - \eta \in \sum_{C \in \Gamma} \text{SOS}_{r_k}(C) \quad (\forall \mathbf{x} \in \mathbb{R}^n), \\ \lambda_1 \in \text{SOS}_{r_1}(C_1), \dots, \lambda_m \in \text{SOS}_{r_m}(C_m).$$

- Lasserre's dense  $\text{SOS}^r$  is obtained if  $\Gamma = \{N\}$ .

-- p.50/58



$$\text{POP: } \zeta^* = \min. f_0(\mathbf{x}) \text{ s.t. } f_k(\mathbf{x}) \geq 0 \ (k = 1, \dots, m)$$

Let  $F_k = \{i : x_i \text{ is involved } f_k\}$ .

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$$r_k = r - \lceil \deg(f_k)/2 \rceil \ (k = 1, \dots, m),$$

-- p.51/58

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$$r_k = r - \lceil \deg(f_k)/2 \rceil \ (k = 1, \dots, m),$$

$\mathbf{u}_q(\mathbf{x}, C) =$  the col. vect. of monomials with  $\deg. \leq q$  in  $x_i$  ( $i \in C$ ).

$$\begin{aligned} \text{Sparse Poly.SDP} &\implies \text{Sparse SDP}^r \text{ by linearization} \\ \min f_0(\mathbf{x}) \text{ s.t. } &\mathbf{u}_{r_j}(\mathbf{x}, C_j) \mathbf{u}_{r_j}(\mathbf{x}, C_j)^T f_k(\mathbf{x}) \succeq \mathbf{O} \ (\forall k) \\ &\mathbf{u}_r(\mathbf{x}, C) \mathbf{u}_r(\mathbf{x}, C)^T \succeq \mathbf{O} \ (\forall C \in \Gamma) \end{aligned}$$

• Lasserre's dense  $\text{SDP}^r$  is obtained if  $\Gamma = \{N\}$ .

-- p.51/58

Example of a sparse POP  $\Rightarrow$  a sparse  $\text{SDP}^2$

$$\begin{aligned} \text{POP: } \min f_0(\mathbf{x}) &= \sum_{i=1}^4 (-x_i^3) \\ \text{s.t. } f_k(\mathbf{x}) &= -a_i \times x_i^2 - x_4^2 + 1 \geq 0 \ (i = 1, 2, 3). \end{aligned}$$

$$\mathbf{R} = \begin{pmatrix} \star & 0 & 0 & \star \\ 0 & \star & 0 & \star \\ 0 & 0 & \star & \star \\ \star & \star & \star & \star \end{pmatrix} \quad \left. \begin{array}{l} C_1 = \{1, 4\}, \ C_2 = \{2, 3\} \\ C_3 = \{3, 4\} \end{array} \right\} \text{max. cliques}$$

$$F_k = \{k, 4\} \subset C_k \ (k = 1, 2, 3).$$

the relax. order  $r = 2 \geq \lceil \max\{\deg(f_k) : 0 \leq k \leq 3\}/2 \rceil = 2$ ,

$$r_k = 2 - \lceil \deg(f_k)/2 \rceil = 1 \ (k = 1, 2, 3),$$

$$\mathbf{u}_1(\mathbf{x}, C_k) = (1, x_k, x_4)^T \ (k = 1, 2, 3),$$

$$\mathbf{u}_2(\mathbf{x}, C_k) = (1, x_k, x_4, x_k^2, x_k x_4, x_4^2)^T \ (k = 1, 2, 3).$$

$$\begin{aligned} \text{Poly.SDP: } \min. f_0(\mathbf{x}) \text{ s.t. } &\mathbf{u}_1(\mathbf{x}, C_k) \mathbf{u}_1(\mathbf{x}, C_k)^T f_k(\mathbf{x}) \succeq \mathbf{O} \ (\forall k) \\ &\mathbf{u}_2(\mathbf{x}, C_k) \mathbf{u}_2(\mathbf{x}, C_k)^T \succeq \mathbf{O} \ (\forall k) \end{aligned}$$

-- p.52/58

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3. Basic idea of SOS and SDP relaxations
4. SOS relaxation of POPs
5. SDP relaxation of POPs
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7. Numerical results

- SparsePOP (Waki et al.) for POPs  $\Rightarrow$  sparse SDPs
- SDPA (Fujisawa et al.) for solving SDPs
- MATLAB Optimization Toolbox to refine the SDP solution
- Intel Xeon 2.66GHz with 12 GB memory

-- p.53/58

Unconstrained optimization problem

The generalized Rosenbrock function — poly. with deg = 4

$$f_R(\mathbf{x}) = 1 + \sum_{i=2}^n (100(x_i - x_{i-1}^2)^2 + (1 - x_i^2)^2)$$

The chained singular function — poly. with deg = 4

$$f_C(\mathbf{x}) = \sum_{i \in J} ((x_i + 10x_{i+1})^2 + 5(x_{i+2} - x_{i+3})^2 + (x_{i+1} - 2x_{i+2})^4 + 10(x_i - 10x_{i+3})^4)$$

Here  $J = \{1, 3, 5, \dots, n - 3\}$ ,  $n$  is a multiple of 4.

minimize  $f_R(\mathbf{x}) + f_C(\mathbf{x})$

— unknown global optimal value and solution

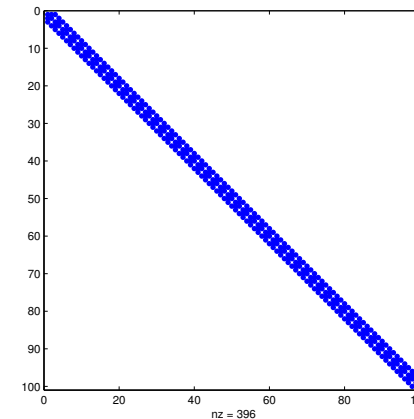
the sparsity pattern matrix  $R$

(=the sparsity pattern of the Hessian matrix)

— very sparse  $\Rightarrow$  next

— p.54/58

Sparsity pattern of the sparsity pattern matrix  $R$  with a simultaneous row and column reordering (Reverse Cuthill-McKee ordering)



Structured sparsity

- Sparse (symbolic) Cholesky factorization

— p.55/58

min  $f_R(\mathbf{x}) + f_C(\mathbf{x})$  — deg. 4, sparse, unknown opt.val.

$n$	Sparse			Dense (Lasserre)		
	$\epsilon_{\text{obj}}$	# =	cpu	$\epsilon_{\text{obj}}$	# =	cpu
12	6e-5	214	0.1	3e-6	1,819	6.9
16	5e-5	294	0.1	1e-9	4,844	71.3
100	7e-6	1,974	0.5	out of	mem	
1000	7e-7	19,974	6.9			
2000	7e-8	39,974	15.1			
3000	out of	mem				

$$\epsilon_{\text{obj}} = \frac{|\text{lbd. for opt.val.} - \text{approx.opt.val.}|}{\max\{1, |\text{lbd. for opt.val.}|\}}$$

# = : the number of equalities of SDP,

cpu : cpu time in second

- Global optimality is guaranteed with high accuracy.

— p.56/58

alkyl from globalib — presented previously

$$\text{min} \quad -6.3x_5x_8 + 5.04x_2 + 0.35x_3 + x_4 + 3.36x_6$$

$$\text{sub.to} \quad -0.820x_2 + x_5 - 0.820x_6 = 0,$$

$$0.98x_4 - x_7(0.01x_5x_{10} + x_4) = 0, \quad -x_2x_9 + 10x_3 + x_6 = 0,$$

$$x_5x_{12} - x_2(1.12 + 0.132x_9 - 0.0067x_9^2) = 0,$$

$$x_8x_{13} - 0.01x_9(1.098 - 0.038x_9) - 0.325x_7 = 0.574,$$

$$x_{10}x_{14} + 22.2x_{11} = 35.82, \quad x_1x_{11} - 3x_8 = -1.33,$$

$$\text{lbd}_i \leq x_i \leq \text{ubd}_i \quad (i = 1, 2, \dots, 14).$$

Sparse			Dense (Lasserre)		
$\epsilon_{\text{obj}}$	$\epsilon_{\text{feas}}$	cpu	$\epsilon_{\text{obj}}$	$\epsilon_{\text{feas}}$	cpu
7.7e-8	1.8e-6	1.4	out of	memory	

$$\epsilon_{\text{obj}} = \frac{|\text{lbd. for opt.val.} - \text{approx.opt.val.}|}{\max\{1, |\text{lbd. for opt.val.}|\}}$$

$\epsilon_{\text{feas}}$  = max. error in equalities, Elapsed time of SDPA in sec.

- Global optimality is guaranteed with high accuracy.

— p.57/58

Some benchmark QOPs from globalib

Problem	n	r	Sparse			Dense
			$\epsilon_{\text{obj}}$	$\epsilon_{\text{feas}}$	E.time	E.time
ex2_1_8	24	2	1.2e-5	1.5e-6	6.7	28.7
ex3_1_1	8	2	3.5e-2	1.8e-14	0.3	0.4
ex3_1_1	8	3	2.5e-7	1.2e-13	0.6	52.3
ex5_4_2	8	2	5.2e-1	2.5e-12	0.3	0.5
ex5_4_2	8	3	8.0e-9	4.6e-16	0.6	49.0
ex5_3_2	22	2	1.5e-1	1.7e-16	1.3	47.4 <sup>†</sup>
ex5_3_2	22	3	1.3e-4	2.6e-14	599.0	-

$r$  = the relax. order, E.time = Elapsed time of SDPA in sec.,

$$\epsilon_{\text{obj}} = \frac{|\text{bd. for opt.val.} - \text{approx.opt.val.}|}{\max\{1, |\text{bd. for opt.val.}|\}}.$$

$\epsilon_{\text{feas}}$  = the max. error in equalities and inequalities,

<sup>†</sup> : a highly accurate sol. is obtained.